

Universality properties of graph homomorphism: one construction to prove them all

Luca Motto Ros

Department of Mathematics “G. Peano”
University of Turin, Italy
luca.mottoros@unito.it
<https://sites.google.com/site/lucamottoros/>

Joint work with S. Scamperti

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Graph homomorphism

\mathcal{G} = all undirected graphs. Notations like $\mathcal{G}^{=\kappa}$, $\mathcal{G}^{<\kappa}$, ... are used for the restrictions of \mathcal{G} to graphs of size = κ , or $< \kappa$, ...

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A homomorphism $h: G_1 \rightarrow G_2$ is a map $h: V(G_1) \rightarrow V(G_2)$ such that

$$v E(G_1) w \Rightarrow h(v) E(G_2) h(w).$$

We write $G_1 \preceq G_2$ if there is such a homomorphism, $G_1 \approx G_2$ if $G_1 \preceq G_2 \preceq G_1$, and $G_1 \prec G_2$ if $G_1 \preceq G_2$ but $G_2 \not\preceq G_1$.

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Empirical fact

The homomorphism relation \preceq on \mathcal{G} is **universal**, i.e. it can code most of the other mathematical relations.

Universality properties of graph homomorphism

Category theory: The category \mathfrak{G} of graphs with arrows given by homomorphisms is *alg-universal*, i.e. every concrete category can be fully embedded into it.

[A full embedding is an injective (on objects) fully faithful functor.]

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Complexity/computer science: Undecidability of the graph homomorphism problem (for finite graphs), etc...

Model theory: The theory of graphs interprets all first-order theories in a countable language.

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We are looking for a very strong yet flexible coding procedure...

The coding procedure

Connected sum

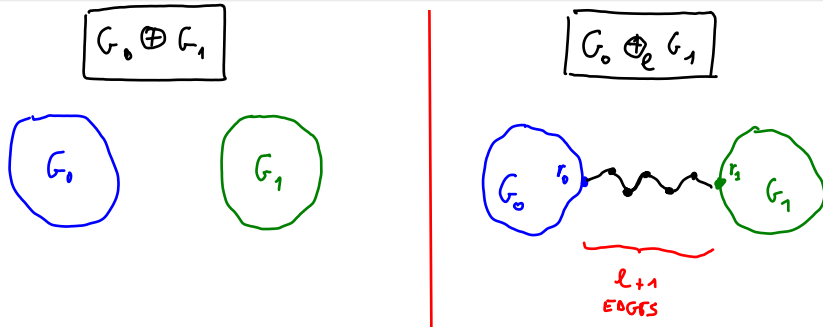
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Definition (binary connected sum)

Given $\ell \in \omega$ and vertices $r_0 \in V(G_0)$ and $r_1 \in V(G_1)$, the (ℓ -)connected sum $G_0 \oplus_\ell G_1$ of G_0 and G_1 is the variant of $G_0 \oplus G_1$ where we add a brand new path of length $\ell + 1$ joining r_0 to r_1 .



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Key idea: generalize this to arbitrary families of graphs.

Definition (generalized connected sum)

$$\{G_a\}_{a \in A} \text{ with } r_a \in V(G_a) \qquad \rho: [A]^2 \rightarrow \omega$$

Then the (ρ -)connected sum $\bigoplus_\rho G_a$ of the G_a 's is the graph obtained by adding to the disjoint sum of the G_a 's a path of length $\rho(\{a, a'\}) + 1$ joining r_a to $r_{a'}$ for all $\{a, a'\} \in [A]^2$ such that $\rho(\{a, a'\}) \neq 0$.

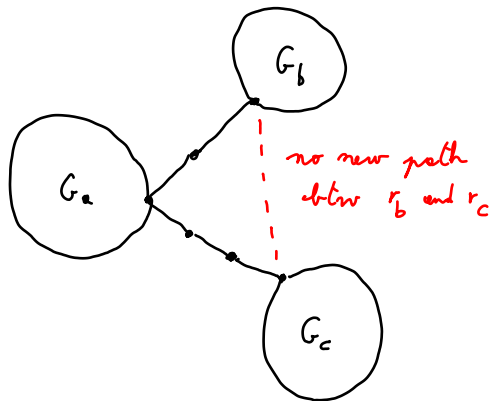
Connected sum (picture)

$$A = \{a, b, c\}$$

$$f(\{a, b\}) = 1$$

$$f(\{a, c\}) = 2$$

$$f(\{b, c\}) = 0$$



Glossary from graph theory

(Finite) Path: Sequence of vertices $(v_i)_{i \leq n}$ with $v_i E(G) v_{i+1}$ for $i < n$

Closed path: Path with $v_n = v_0$

Cycle: Closed path with $v_i \neq v_j$ for $i \neq j$ (except for 1st and last element)

Odd closed path/cycle: Closed path/cycle with odd length

Bipartite: Graph without odd cycles (equiv.: 2-chromatic)

Rigid: Graph whose only endomorphism is the identity

Distance: $d(v, w) =$ length of the shortest path joining v and w

Diameter: $\text{diam}(G) = \sup\{d(v, w) \mid v, w \in V(G)\}$

Connected: Graphs in which all pairs of vertices are joined by some path

Definition

A graph G is **uniformly non-bipartite** if there is $n \in \omega$ such that every $v \in V(G)$ belongs to an odd closed path of length $\leq n$; the smallest such n , which is odd and at least 3, is denoted by n_G . We set

$$\delta(G) = \min\{n_G - 2, \text{diam}(G)\}$$

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Remarks:

- 1 If G has finite diameter (e.g. G connected and finite), then this is equivalent to being non-bipartite.
- 2 Every non-bipartite finite graph is homomorphically equivalent to a uniformly non-bipartite graph.
- 3 For infinite graphs, this is more general than having finite diameter, yet there are non-bipartite graphs which are not homomorphically equivalent to a uniformly non-bipartite one.

Characteristic of graphs

Girth: $\gamma(G)$ = length of the shortest cycle

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Key fact

If $G \preceq H$ then $\chi(G) \leq \chi(H)$ and $\gamma_o(G) \geq \gamma_o(H)$.

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Key fact

If $G \preceq H$ then $\chi(G) \leq \chi(H)$ and $\gamma_o(G) \geq \gamma_o(H)$.

Thus if $\chi(G) < \chi(H)$ and $\gamma_o(G) < \gamma_o(H)$,
then G and H are \preceq -incomparable.

Connected sum (reprise)

Let I be a cardinal and $\mathcal{H} = \{H_i\}_{i \in I}$ be a family of connected uniformly non-bipartite pairwise \preceq -incomparable graphs with $r_i \in V(H_i)$.

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Let A be a nonempty set, and let $i: A \rightarrow I$ and $\rho: [A]^2 \rightarrow \omega$ be arbitrary functions. Then

$$\bigoplus_{\substack{A, \mathcal{H} \\ \rho, i}} G_a$$

is the ρ -connected sum of $\{G_a\}_{a \in A}$, where each G_a is a (distinct) isomorphic copy of $H_{i(a)}$.

Theorem (M.-Scamperti)

Suppose that ρ takes large enough values compared to the $\delta(H_i)$'s.

(E.g. $\rho(\{a, a'\}) \geq \delta(\mathcal{H})$ whenever $\rho(\{a, a'\}) \neq 0$, if (\star_2) holds.)

Then for every $j \in I$ and every homomorphism $h: H_j \rightarrow \bigoplus_{\rho, i}^{A, \mathcal{H}} G_a$ there is a unique $\bar{a} \in A$ such that $h(H_j) \cap G_{\bar{a}} \neq \emptyset$, and $i(\bar{a}) = j$.

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Remark: It is known that there are enough \preceq -incomparable rigid graphs to start with...

How to code a first-order structure

A variation of the interpretation procedure from model theory, shows that wlog we can assume that the structure is an **I -colored graph** $C = (G, c)$, i.e. a graph G together with a singled-out I -coloring $c: V(G) \rightarrow I$.

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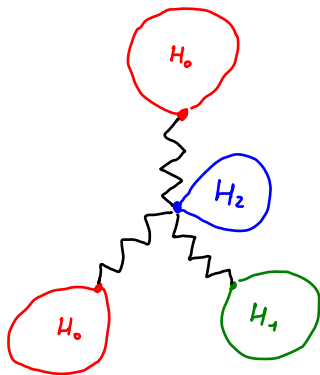
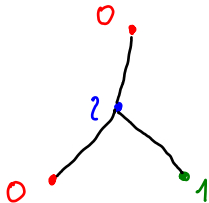
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How to code a first-order structure (picture)

$$\mathcal{H} = \{H_0, H_1, H_2\}$$



$$C = (G, c)$$

$$c: V(G) \rightarrow \{0, 1, 2\}$$

||
A

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Interesting fact: Once the full machinery is developed, for most of the applications it is enough to deal with the case $I = 3$.

Applications

Theorem (Erdős)

For every finite chromatic number n there are (finite) graphs G with $\chi(G) = n$ and arbitrarily high girth.

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- If G is uncountably chromatic, then either $\gamma(G) = \gamma_o(G) = 3$, or else $\gamma(G) = 4 < \gamma_o(G)$.
- For every $n > \aleph_0$, there are graphs G (of size n) with $\chi(G) = n$ and arbitrarily high odd girth.

Graphs with prescribed characteristics

More generally, fix a cardinal n and $m, k \in \omega \cup \{\infty\}$: we want to deal with the class $\mathcal{G}_{n,m,k}$ of graphs G with $\chi(G) = n$, $\gamma(G) = m$, and $\gamma_o(G) = k$.

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Questions

How many graphs are there in such classes? How complicated is their homomorphism structure? (E.g. large antichains, descending chains, and so on.) Do they contain rigid graphs?

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The answer was known for n finite and $m = k = 3$, or $n = 3$ and arbitrarily high $m = k$: indeed, the corresponding categories are alg-universal (Pultr-Trnková).

Theorem (M.-Scamperti)

Let (n, m, k) be an acceptable triple. Then either $\mathcal{G}_{n,m,k}$ is trivial (if $n = 2$ or $n = m = k = 3$), or else the homomorphism category associated to it is alg-universal.

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- 4 If $n < \aleph_0$, then $\mathcal{G}_{n,m,k}^{<\aleph_0}$ is order-theoretically universal.

Graphs with prescribed characteristics

The proof relies on the following results.

Lemma (M.-Scamperti)

Let $G = \bigoplus_{\rho,i}^{A,\mathcal{H}} G_a$ with $\rho(\{a, a'\}) \geq \delta(\mathcal{H})$ when different from 0. Then

$$\chi(G) = \sup_{i \in I} \chi(H_i) \qquad \gamma(G) = \min_{i \in I} \gamma(H_i) \qquad \gamma_o(G) = \min_{i \in I} \gamma_o(H_i).$$

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Let κ be a cardinal such that $\kappa^{<\kappa} = \kappa$. For all (n, m, k) as above and $n \leq \kappa$, the relation $\preceq \upharpoonright \mathcal{G}_{n,m,k}^{\kappa}$ is complete (in fact, invariantly universal) for κ -analytic quasi-orders with respect to κ -Borel reducibility.

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- The chromatic number can be replaced by the **circular chromatic number** χ_c , the **fractional chromatic number** χ_f , and so on.

Forbidden graphs

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One of the best known results concerning the structure of $\text{Forb}_{\mathcal{F}}$ was:

Theorem (Nešetřil-Rödl)

Every $\text{Forb}_{\mathcal{F}}$, if not trivial, contains an infinite set of \preceq -incomparable (finite) graphs.

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- 3 The homomorphism category associated to $\text{Forb}_{\mathcal{F}}$ is alg-universal, as witnessed by a cardinal-respecting full embedding.

(All remaining cases.)

Uncountably chromatic graphs

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In contrast:

Theorem (Hajnal-Komjáth)

Assuming the Continuum Hypothesis CH, there is a triangle-free graph (of size 2^{\aleph_0}) having chromatic number \aleph_1 and omitting the graph $H_{\omega, \omega+2}$.

Uncountably chromatic graphs

It took 30 years to even just prove that the hypothesis CH can be removed.

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In particular, there are $2^{2^{\aleph_0}}$ -many pairwise \preceq -incomparable (even rigid, connected) graphs of this kind.

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Moreover, we can control girth and odd girth (besides obvious limitations).

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This generalizes (with caution) to infinite graphs as well.

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- 2 If there are three \preceq -incomparable graphs in $\mathcal{C}^{<\aleph_0}$, then such embedding respect cardinalities. Hence $\mathcal{C}^{<\aleph_0}$ is order-theoretically universal, $\mathcal{C}^{=\aleph_0}$ is complete for analytic quasi-orders, etc...
- 3 If there are three *rigid* \preceq -incomparable graphs in \mathcal{C} , then \mathcal{C} is even alg-universal.

A final comment

In all applications, our method reveals a sort of general dichotomy:

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Thank you for your attention!